# Phylogenetic Inference: Building Trees 

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1195 env sequences from 9 HIV+ patients [taken from Rambaut et al. (2004)]

Retroviruses (and HBV) exist as a quasi-species within infected patients:

- Shared substitutions may be insufficient to resolve intra-host phylogenies

Improve resolution using joint model:

- Indel rates $\geq$ substitution rates
- Opportunity to detect intra-host recombination

- Contentious issue among paleobiologists: Do Archaea (Euryarchaeota/Eocytes) form one or two domains? Weekly World News calls humans slime molds.


## The Chicken or the (Small) Genome: Which Came First?



## Maximum Parsimony (MP)

Most often used $\neq$ "best", not even statistically consistent, but fast, fast, fast . . . if you know the tree

Key: Find tree with minimal \# of "suspected" substitutions (internal states are not observed, 0/1 model process)

- Counting minimum \# of substitutions is easy
- Enumerating (searching through) all possible trees is hard

Human-T T C C TGGAAT
Chimp -T ACCTGGAAT
Mouse - A ACCT- T T A T
Fly $\mathbf{- A} \mathbf{A} \mathbf{A}$ T C GTAT
Site:1 $2 \begin{array}{lllllllll} & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \\ \text { Along Molecular Sequence }\end{array}$
Sites are independent


SISMID - p. 5

## Maximum Parsimony (MP)

## A little history:

- Anthony Edwards/Luca Cavalli-Sforza $(1963,1964)$
- Both students of R.A Fisher
- Introduced both parsimony and likelihood methods (for continuous quantities, e.g. gene frequencies) in one paper
- Camin and Sokal (1965) provide first program for molecular sequences
- Fitch and Margoliash (1967) provide efficient algorithm



## Maximum Parsimony Algorithm

procedure Fitch and Margoliash (1967) Algorithm
cost $C \leftarrow 0$ \{Initialization\}
pointer $k \leftarrow 2 N-1$ \{at the root node\}
To obtain the set $R_{k}$ of possible states at node $k$ \{Recursion\}
if $k$ is leaf then
$R_{k} \leftarrow$ observed character for taxon $k$
else
Compute $R_{i}, R_{j}$ for daughters $i, j$ of $k$
if $R_{i} \cap R_{j} \neq 0$ then
$R_{k} \leftarrow R_{i} \cap R_{j}$
else
$R_{k} \leftarrow R_{i} \cup R_{j}$
$C \leftarrow C+1$
end if
end if
minimum cost is $C$ \{Termination\}


## Searching for the MP Tree

Complexity:

- Find MP score is NP-complete
- Find MP tree is NP-hard


Recall that \# of $N$-taxon rooted trees is $3 \times 5 \times \cdots \times 2 N-3$ Attack exponential-order space Branch-and-Bound:

- Monotonic order: min $\mathrm{PS}_{2} \leq \min \mathrm{PS}_{3} \leq \ldots$
- Bound if min $\mathrm{PS}_{k}>$ best $n$-taxon PS found so far.


## Neighbor-Joining (Saitou and Nei, 1987)

Computational algorithm: alignment $\rightarrow$ single tree


- Advantages: very fast, great for 1000 s of sequences
- Disadvantages: no site-to-site rate variation, no natural ways to compare trees/measure data support


## Neighbor-Joining

Caveat: Pairs $i, j$ with $\min d_{i j}$ are not necessarily nearest neighbors.
E.g., $d_{\mathrm{AB}}=3<d_{\mathrm{AC}}=5$


Solution: Subtract off the average distances to all other leaves via

$$
D_{i j}=d_{i j}-\left(r_{i}+r_{j}\right), \quad r_{i}=\frac{1}{|L|-2} \sum_{k \in L} d_{i k},
$$

where $L$ is the current set of leaves. Proof in Studier and Keppler (1988).

Computational: $O\left(N^{3}\right)$

## Likelihood-based Methods (Felsenstein, 1973)

Statistical technique: assumes an unknown tree and a stochastic model for character change along the tree


- Advantages: site-to-site rate/tree variation is easy, can formulate probability statements
- Disadvantages: must "search" tree-space $\rightarrow$ slow

Foundation of Bayesian Phylogenetics

Reconstruction Example
Human-T TC C T G G A A T Chimp - T A C C T G G A A T Mouse - A A C T - - T A T Fly - A G A T C G T A T Site: $\begin{array}{llllllllll}1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \\ & & \text { Along Molecular Sequence }\end{array}$


- Substitution: single residue replaces another
- Insertion/deletion: residues are inserted or deleted


## Statistical Model

Assume: Homologous sites are iid and site patterns (e.g. dotted box)
$X Y \ldots Z \sim$ Multinomial $\left(p_{X Y \ldots Z}\right)$
where $p_{X Y \ldots Z}$ is determined by an unknown tree $\tau$, branch lengths $t \in \mathbf{T}$ and continuous-time Markov chain model (for residue substitution) given by infinitesimal rate matrix $\mathbf{Q}$

$$
\mathbf{P}(X \rightarrow Y \text { in time } t)=e^{t \mathbf{Q}}
$$

## CTMC( $Q$ ) = $\epsilon \sim \operatorname{Normal}\left(\mu, \sigma^{2}\right)$ of Phylogenetics

Continuous in elapsed time $t$, discrete in starting/ending state!
Memory-less process in which the probability that state $b$ replaces state $a$ during ( $t, t+$
 $s)=s q_{a b}+o(s)$

- Infinitesimal generator matrix $Q$ has off-diagonal entries $q_{a b}$ and row sums $=0$

Think: Exponential waiting time with rate $R_{a}=\sum_{b} q_{a b}$ until chain leaves $a$. Then the new state $b$ is independently chosen with probabilities $q_{a b} / R_{a}$

## From Infinitesimal to Finite Time

Let $p_{a b}(t)=$ the finite-time probability of the chain moving from state $a$ at time 0 to state $b$ at time $t$, then matrix $\boldsymbol{P}(t)=\left\{p_{a b}(t)\right\}$ satisfies

$$
\frac{\mathrm{d}}{\mathrm{~d} t} \boldsymbol{P}(t)=\boldsymbol{P}(t) \boldsymbol{Q} \text { where } \boldsymbol{P}(0)=\boldsymbol{I}
$$

with solution

$$
\boldsymbol{P}(t)=e^{t \boldsymbol{Q}}=\boldsymbol{I}+t \boldsymbol{Q}+\frac{1}{2}(t \boldsymbol{Q})^{2}+\cdots=\sum_{k=0}^{\infty} \frac{1}{k!}(t \boldsymbol{Q})^{k}
$$

as

$$
\frac{\mathbf{d}}{\mathbf{d} t} e^{t \boldsymbol{Q}}=\boldsymbol{Q} e^{t \boldsymbol{Q}}=e^{t \boldsymbol{Q}} \boldsymbol{Q} \text { for } t \text { real }
$$

## Example: Two-State Model

Consider purines (R) $\leftrightarrow$ pyrimidines (Y). Kolmogorov forward equation:

$$
p_{\mathrm{RY}}(t+s)=p_{\mathrm{RR}}(t) \alpha s+p_{\mathrm{RY}}(t)(1-\beta s)+o(s)
$$

yielding

$$
\frac{\mathrm{d}}{\mathrm{~d} t} p_{\mathrm{RY}}(t)=\alpha p_{\mathrm{RR}}(t)-\beta p_{\mathrm{RY}}(t)
$$

$$
\boldsymbol{Q}=\left(\begin{array}{cc}
-\alpha & \alpha \\
\beta & -\beta
\end{array}\right)
$$

with eigenvalues 0 and $-(\alpha+\beta)$

Solutions of $\boldsymbol{P}(t)=e^{t \boldsymbol{Q}}$ have the form

$$
c+d e^{-(\alpha+\beta) t}
$$

## Standard CTMCs for Phylogenetics

- Jukes and Cantor (JC69),

$$
\pi_{a}=\frac{1}{4}, \kappa_{1}=\kappa_{2}=1
$$

- Kimura (K80), $\pi_{a}=\frac{1}{4}, \kappa_{1}=\kappa_{2}$
- Hasegawa, Kishino and Yano (HKY85), $\kappa_{1}=\kappa_{2}$ (most common)
- Tamura and Nei (TN93), right

- General Time Reversible (GTR)

Note identifiability concern in $e^{t Q}$. Common solution is to fix 1 d.f. such that

$$
\sum_{a} q_{a a} \pi_{a}=-1
$$

Scaling: $t=1 \Rightarrow 1$ expected substitution per site

## Explicit Parameterization of TN93

Nucleotides mutate according to a Markovian process

$$
\operatorname{Pr}(X \rightarrow Y \text { in time } t)=e^{t Q_{\text {Nuc }}}
$$

where $Q_{\text {Nuc }}$ is a $4 \times 4$ infinitesimal rate matrix and $t$ is a branch length.

$\kappa_{1}, \kappa_{2}$ are transition:transversion rate ratios and $\pi$ is the stationary distribution of $\{\mathrm{A}, \mathrm{G}, \mathrm{C}, \mathrm{T}\}$. $\beta$ controls the overall rate and can vary from site-to-site.

## Site-to-Site Rate Variation

Variation occurs quite naturally and is also an important inference

- short range: codon phase (slow-slow-fast)
- long range: enzymatic active sites, protein folding, immunological pressures/selection


Assume: infinitesimal rates for site $k$ are $r_{k} \times t \times q_{a b}$. Various priors on $r_{k}$ with $\mathbf{E}\left(r_{k}\right)=1$. Implicitly Bayesian

- Yang (1994) - discretized Gamma distribution


## General Time Reversible CTMC

## Let

$$
Q=R D_{\pi}
$$

where $R$ is symmetric and $D_{\pi}$ is a diagonal matrix composed of the stationary distribution $\pi$.

- Detailed balance $\Leftrightarrow \pi_{a} q_{a b}=\pi_{b} q_{b a}$. Balance + irreducibility $\Leftrightarrow$ reversible
- Note $\boldsymbol{Q}$ is similar to $\boldsymbol{R}$, as $\boldsymbol{D}^{1 / 2} \boldsymbol{Q} \boldsymbol{D}^{-1 / 2}=\boldsymbol{R}$
- Hence, $Q$ must have real eigenvalues and real eigenvectors

The properties speed up computation of the finite-time transition matrix $\boldsymbol{P}(t)=e^{t \boldsymbol{Q}}$

## Calculating the Probability of a Single Site Pattern $Y_{i}$

Given the tree and unobserved internal node states, the probability is the product of the finite time mutation probabilities over all branches:


$$
\begin{align*}
& L\left(\boldsymbol{Y}_{i}\right) \propto p_{\text {AAGT }}=\sum_{X} \sum_{Y} \operatorname{Pr}\left(Y \rightarrow \mathbf{A}, t_{1}\right) \operatorname{Pr}\left(X \rightarrow \mathbf{G}, t_{2}\right) * \\
& \operatorname{Pr}\left(X \rightarrow \mathbf{T}, t_{3}\right) \operatorname{Pr}\left(Y \rightarrow \mathbf{A}, t_{4}\right) \operatorname{Pr}\left(X \rightarrow Y, t_{5}\right) \pi_{X} \tag{1}
\end{align*}
$$

- Number of sumants grow rapidly in $N \rightarrow$ sum-product/peeling algorithm to distribute sums across the product


## Pruning Algorithm Felsenstein (1981)

Let $P\left(L_{k} \mid a\right)=$ likelihood of leaves below node $k$ given $k$ is in state $a$. Then, recursively compute $P\left(L_{k} \mid a\right)$ given $P\left(L_{i} \mid b\right)$ and $P\left(L_{j} \mid c\right)$ for daughters $i, j$ of $k$ :

```
Set pointer \(k \leftarrow 2 N-1\) \{the root, initialization\}
Compute \(P\left(L_{k} \mid a\right) \forall a\) as follows: \{recursion\}
if \(k\) is a leaf node then
    if \(a\) is observed then
                \(P\left(L_{k} \mid a\right)=1\)
            else
                \(P\left(L_{k} \mid a\right)=0\)
            end if
```



```
else
            Compute \(P\left(L_{i} \mid a\right)\) and \(P\left(L_{j} \mid a\right) \forall a\) for daughters \(i, j\) of \(k\) \{post-order traversal\}
            \(P\left(L_{k} \mid a\right)=\sum_{b} \sum_{c} \operatorname{Pr}\left(a \rightarrow b, t_{i}\right) P\left(L_{i} \mid b\right) \times \operatorname{Pr}\left(a \rightarrow c, t_{j}\right) P\left(L_{j} \mid c\right)\)
end if
\(L\left(\boldsymbol{Y}_{i}\right) \leftarrow \sum_{a} P\left(L_{2 n-1} \mid a\right) \pi_{a}\{\) termination \(\}\)
```


## ML Tree or MAP Tree?

Reporting uncertainty on tree estimates:

- The Bootstrap
- Most common
- Assumes evolutionary events are reproducible. "If I went back out to the field and recollected exchangeable data ..."
- Bayesian inference
- Returns the probability of a tree given the observed data and model
- Requires MCMC (e.g., MrBayes or BEAST)
- Advantages
* Does not rely on asymptotics (hypothesis testing)
* Naturally incorporates uncertainty in all parameters (including discrete quantities: trees, site-classifications, etc.)
* Arguably faster algorithms
- Disadvantages
* Must specify (justifiable) prior distributions

